

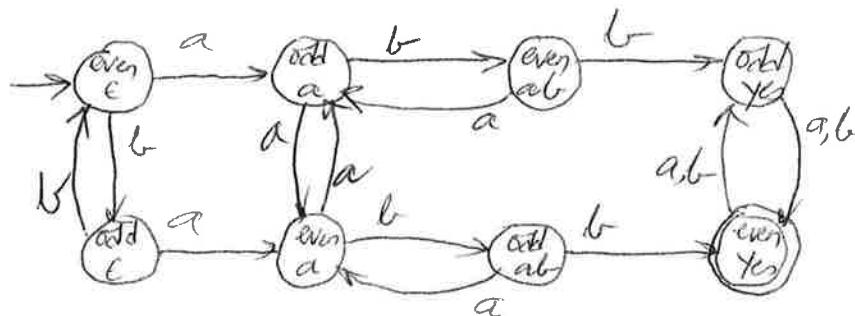
600.271 Automata & Computation Theory
Mid-Semester Examination
March 8, 2012
In-class, Closed Book, Time: 1 hr, 10 mins

All the subproblems carry equal weight. There are 5 subproblems in this examination.

I. Design the specified automaton for every one of the following languages.

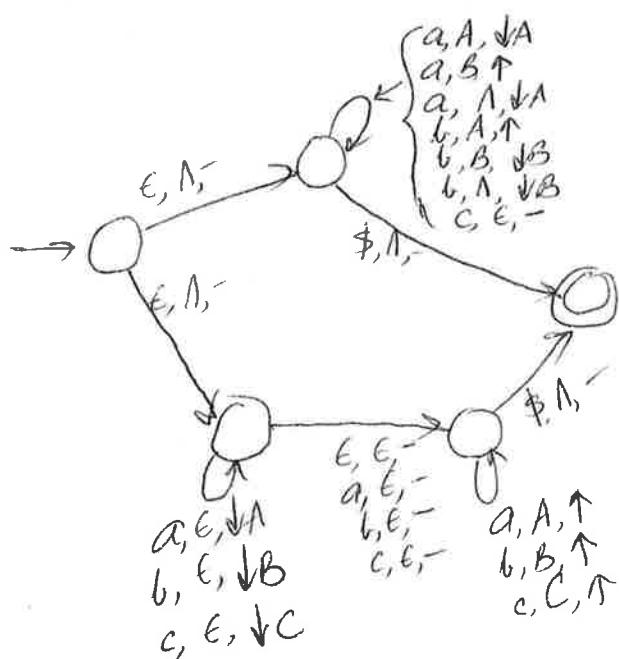
1. A dfa for the language $L_1 = \{x \mid x \in \{a, b\}^*, \text{ abb is a substring of } x, \text{ and } |x| \text{ is an even integer}\}$.

search for abb; also keep track of the parity of $|x|$.



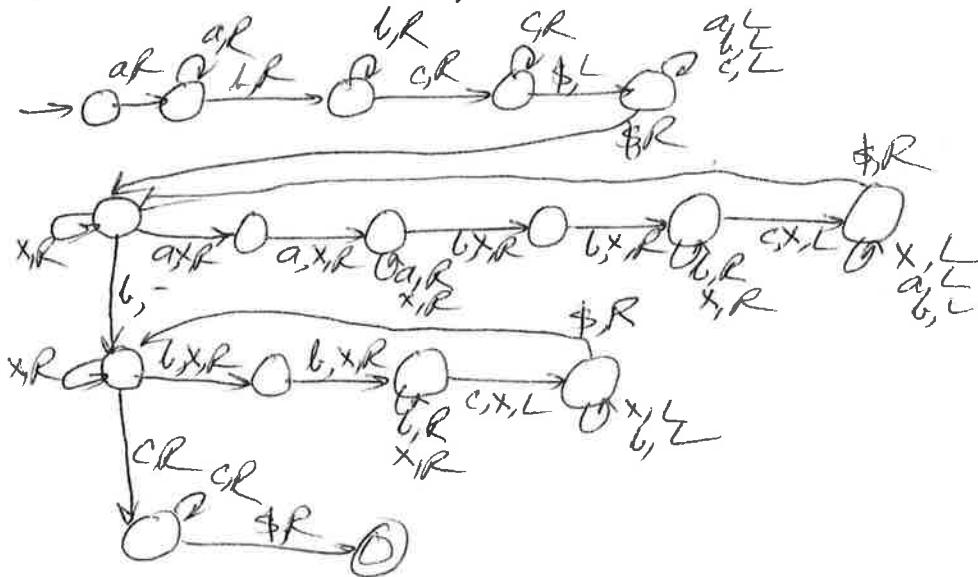
2. An npda for the language $L_2 = \{x \mid x \in \{a, b, c\}^*, \text{ and } (\#_a x = \#_b x, \text{ or } x = x^R)\}$.

Guess the condition satisfied and verify. For $x = x^R$, guess the midpoint.



3. A dLBA for the language $L_3 = \{a^i b^j c^k \mid i, j, k \geq 1, i \text{ and } j \text{ are even, and } i < j < 2k\}$.

check if $a^i b^j c^k$. cross off 2a's, 2b's, 1c repeatedly until a's are exhausted. Then cross off 2b's, 1c repeatedly until b's are exhausted.



4. A CFG for the language $L_4 = \{a^i b^j a^k b^\ell \mid i, j, k, \ell \geq 1, \text{ and } (i = j \text{ and } k \neq \ell) \text{ or } (i = \ell \text{ and } j \neq k\}$.

From X generate $\{a^i b^i \mid i \geq 1\}$, and from Y generate $\{a^i b^\ell \mid i, j \geq 1, i \neq j\}$.

$$X \rightarrow aXb \mid ab$$

$$\begin{cases} Y \rightarrow U \mid V \\ U \rightarrow aU \mid ax \\ V \rightarrow vb \mid xb \end{cases}$$

$$S \rightarrow XY \mid W \quad \left\{ \begin{array}{l} i=j \text{ and } k \neq \ell ; \\ i=\ell, j \neq k \end{array} \right\} \text{ from } XY \quad W$$

$$W \rightarrow a^k b^\ell \mid a^i y^i$$

II. Prove that the following language is not an fa language by applying the pumping lemma.

$$L_5 = \{a^i b^j \mid i, j \geq 1, i \geq j \text{ or } 2i \leq j\}.$$

For any k , let $z = a^{3k} b^{3k}$. Note that $z \in L_5$ since $3k \geq 3k$. Also $|z|=3k$.

For any way of writing $a^{3k} b^{3k} = uvw$ s.t. $|uv| \leq k$ & $|v| > 0$:
since $|uv| \leq k$, $uvw \in a^*$ & hence $v \in a^*$. Let $v = a^l$, $1 \leq l \leq k$.

Then $uv^0 w = a^{3k-l} b^{3k} \notin L_5$.

Hence, by the contrapositive form of the PL, L_5 is not an
fa language.